

$$\frac{}{\vec{\Gamma}, A \vdash A} (\alpha)$$

$$\frac{}{\vec{u} \Vdash \vec{\Gamma}, v \Vdash A \vdash v \Vdash A} (\alpha)$$

idea: (α) y (ω)

$$\frac{\vec{\Gamma}, A, A \vdash B}{\vec{\Gamma}, A \vdash B} (\beta)$$

$$\xrightarrow{(*)} \frac{\vec{u} \Vdash \vec{\Gamma} \quad v \Vdash A \quad w \Vdash A \vdash e \Vdash B}{\vec{u} \Vdash \vec{\Gamma}, v \Vdash A \quad w \Vdash A \vdash e \Vdash B} (\gamma)$$

sin \models

hip de inducción

Però, sustituyendo $\{w/v\}$ en toda la prueba de $(*)$ se llega a una prueba de $\vec{u} \Vdash \vec{\Gamma}, w \Vdash A, w \Vdash A \vdash e \Vdash B$

$$(\gamma) \xrightarrow{\vec{u} \Vdash \vec{\Gamma}, w \Vdash A \vdash e \Vdash \{w/v\} \vdash B}$$

↓
hip de inv

$$\frac{\vec{\Gamma} \vdash A_1 \quad \vec{\Gamma} \vdash A_2}{\vec{\Gamma} \vdash A_1 \wedge A_2}$$

$$\frac{\vec{u} \Vdash \vec{\Gamma} \vdash e_1 \Vdash A_1 \quad \vec{v} \Vdash \vec{\Gamma} \vdash e_2 \Vdash A_2}{\vec{u} \Vdash \vec{\Gamma} \vdash e \Vdash A_1 \wedge e \Vdash A_2}$$

$$e := \langle e_1, e_2 \rangle$$

↓ def.

$$e = e(\vec{u}, \vec{v}, F(\vec{\Gamma}, A)) \quad e \Vdash A_1 \wedge A_2$$

$$\frac{\vec{u} \Vdash \vec{\Gamma}, v \Vdash A \vdash e \Vdash B}{\vec{u} \Vdash \vec{\Gamma}, v \Vdash A \rightarrow e \Vdash B} (\rightarrow_I)$$

$$(\rightarrow_I) \frac{\vec{u} \Vdash \vec{\Gamma} \vdash v \Vdash A \rightarrow e \Vdash B}{\vec{u} \Vdash \vec{\Gamma} \vdash \forall v. (v \Vdash A \rightarrow e \Vdash B)}$$

$$e \{v/n\} = e$$

$$\Gamma(\vec{\gamma}_v, e, v, k) \wedge e = U(k)$$

$$\underbrace{\Gamma_{\infty, 0}}_{\perp} \vdash \underbrace{s(x) = 0}_{\perp} \rightarrow \perp$$

no tiene
realizadores.

$$\emptyset \rightarrow \emptyset$$

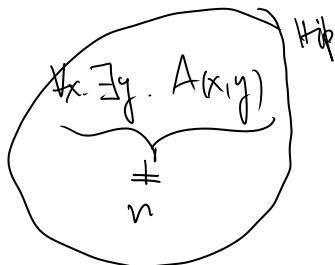
$$\vdash x. \left(\underbrace{x \vdash s(x) = 0}_{\perp} \rightarrow \{ \Gamma_{\infty, 0} \}(x) \vdash \perp \right)$$

$$A(x, y) := (y = 0 \wedge \exists k. T(x, x, k)) \vee (y \neq 0 \wedge \forall k. \neg T(x, x, k))$$

En PA vale $\vdash x \exists y A(x, y)$ pero no vale que

exista e tal que $\vdash x. A(x, \{e\}(x))$ porque si no
 $K = \{x \mid \{x\}(x) \downarrow\}$ sería decidible.

(CT_0) es realizable en Heyting:



$$\vdash x. \{n\}(x) \downarrow \vdash \exists y. A(x, y)$$

$$\vdash x. \{n\}(x) \downarrow \wedge (\{n\}(x))^2 \vdash A(x, \{n\}(x))$$

$$r \vdash \exists e. \vdash x. A(x, \{e\}(x)) \Leftrightarrow r^2 \vdash \vdash x. A(x, \{r^1\}(x))$$

$$\Leftrightarrow \vdash x. \{r^2\}(x) \downarrow \vdash A(x, \{r^1\}(x))$$

$$\boxed{\langle \Gamma_{\infty, \{r^1\}(x)}, \Gamma_{\infty, \{r^2\}(x)} \rangle} \vdash \exists e. \vdash x. A(x, \{e\}(x)) \text{ tesis.}$$

Però PA $\not\vdash CT_0$ porque el halting problem es undecidable

Así que HA $\not\vdash CT_0$

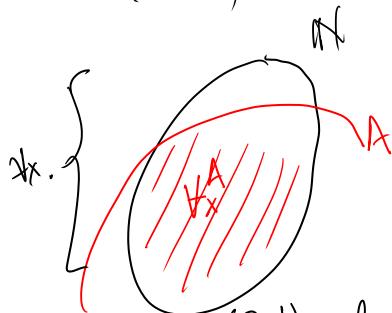
Fórmulas Quasi-Negativas

$$A, B ::= \perp \mid \underbrace{t=s \mid \exists x. t=s \mid A \wedge B \mid A \rightarrow B \mid \forall x. A}_{\text{atômico}}$$

$$\text{ECT}_o: \vdash^A \exists y. B(x,y) \rightarrow \exists e. \vdash^A B(x, \epsilon(x))$$

A es Quasi-negativa

$$\forall_x^A B := \forall x. (A(x) \rightarrow B)$$



Lema 9.1: Para fórmulas A Quasi-Negativas (QN) vale.

$$(1) HA \vdash (\exists x. x \neq A) \rightarrow A \quad \text{definido a partir de } A$$

$$(2) \text{Existem fórmulas } \not\models_A \text{ tal que } HA \vdash A \rightarrow \not\models_A \neq A$$

$$\text{Em part.: } HA \vdash (\exists x. x \neq A) \rightarrow \not\models_A \neq A \\ \vdash (\exists x. x \neq A) \leftrightarrow A$$

Dai: